AI I: problem-solving and search

Lecturer: Tom Lenaerts

Institut de Recherches Interdisciplinaires
et de Développements en Intelligence
Artificielle (IRIDIA)

Université Libre de Bruxelles

Outline

- **Problem-solving agents**
  - ☐ **A kind of goal-based agent**
- **Problem types**
  - ☐ **Single state (fully observable)**
  - ☐ **Search with partial information**
- **Problem formulation**
  - ☐ **Example problems**
- **Basic search algorithms**
  - ☐ **Uninformed**

Problem-solving agent

- Four general steps in problem solving:
  - **Goal formulation**
    - What are the successful world states
  - **Problem formulation**
    - What actions and states to consider give the goal
  - **Search**
    - Determine the possible sequence of actions that lead to the states of known values and then choosing the best sequence.
  - **Execute**
    - Give the solution perform the actions.

Problem-solving agent

function SIMPLE-PROBLEM-SOLVING-AGENT(*percept*) **return** an action
    static: *seq*, an action sequence
        state, some description of the current world state
        goal, a goal
        problem, a problem formulation

    state ← UPDATE-STATE(*state*, *percept*)
    if *seq* is empty **then**
        goal ← FORMULATE-GOAL(*state*)
        problem ← FORMULATE-PROBLEM(*state*,*goal*)
        seq ← SEARCH(*problem*)
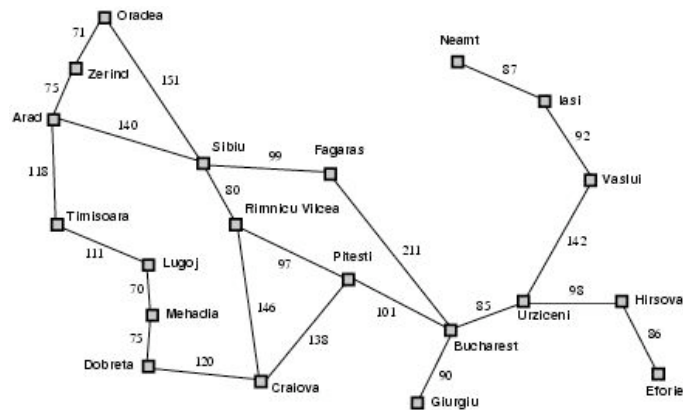   *action* ← FIRST(*seq*)
    seq ← REST(*seq*)
    return *action*

Example: Romania

Example: Romania

- On holiday in Romania; currently in Arad
  - **Flight leaves tomorrow from Bucharest**
- Formulate goal
  - **Be in Bucharest**
- Formulate problem
  - **States: various cities**
  - **Actions: drive between cities**
- Find solution
  - **Sequence of cities; e.g. Arad, Sibiu, Fagaras, Bucharest, ...**

Problem types

- Deterministic, fully observable ⇒ *single state problem*
  - **Agent knows exactly which state it will be in; solution is a sequence.**
- Partial knowledge of states and actions:
  - **Non-observable ⇒ *sensorless or conformant problem***
    - Agent may have no idea where it is; solution (if any) is a sequence.
  - **Nondeterministic and/or partially observable ⇒ *contingency problem***
    - Percepts provide *new* information about current state; solution is a tree or policy; often interleave search and execution.
  - **Unknown state space ⇒ *exploration problem* ("online")**
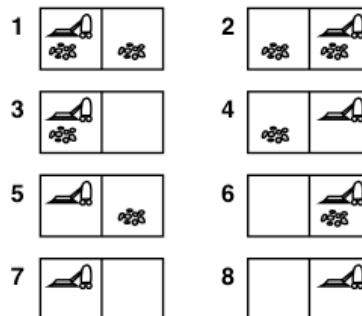   - When states and actions of the environment are unknown.

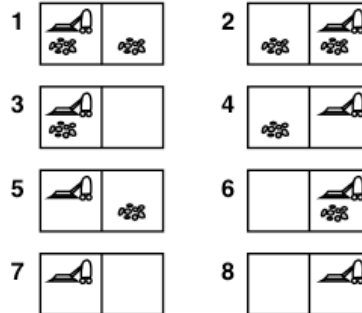# Example: vacuum world

- Single state, start in #5. Solution??

Example: vacuum world

- Single state, start in #5. Solution??
  - *[Right, Suck]*

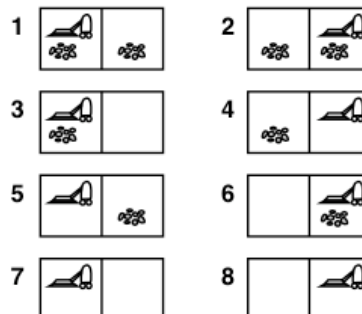# Example: vacuum world

- Single state, start in #5. Solution??
  - *[Right, Suck]*
- Sensorless: start in {1,2,3,4,5,6,7,8} e.g Right goes to {2,4,6,8}. Solution??
- Contingency: start in {1,3}. (assume Murphy's law, Suck can dirty a clean carpet and local sensing: [location,dirt] only. Solution??

Problem formulation

- A problem is defined by:
    - An **initial state**, e.g. *Arad*
    - **Successor function** *S(X)*= set of action-state pairs
        - e.g. *S(Arad)={<Arad → Zerind, Zerind>,…}*
    - **intial state + successor function = state space**
    - **Goal test**, can be
        - Explicit, e.g. *x='at bucharest'*
        - Implicit, e.g. *checkmate(x)*
    - **Path cost** (additive)
        - e.g. sum of distances, number of actions executed, …
        - *c(x,a,y)* is the step cost, assumed to be >= 0

    A solution is a sequence of actions from initial to goal state.
    Optimal solution has the lowest path cost.
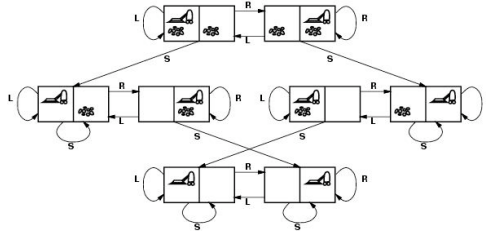

Selecting a state space

- Real world is absurdly complex.
    - **State space must be *abstracted* for problem solving.**
- (Abstract) state = set of real states.
- (Abstract) action = complex combination of real actions.
    - **e.g. Arad →Zerind represents a complex set of possible routes, detours, rest stops, etc.**
    - **The abstraction is valid if the path between two states is reflected in the real world.**
- (Abstract) solution = set of real paths that are solutions in the real world.
- Each abstract action should be "easier" than the real problem.

Example: vacuum world

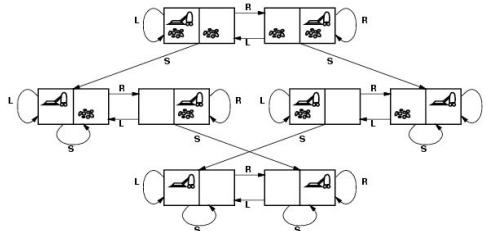

- States??
- Initial state??
- Actions??
- Goal test??
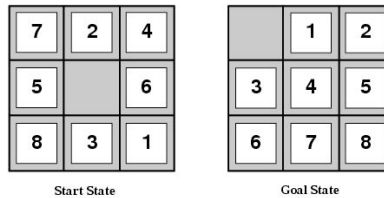- Path cost??

Example: vacuum world

- States?? two locations with or without dirt: $2 \times 2^2 = 8$ states.
- Initial state?? Any state can be initial
- Actions?? {*Left*, *Right*, *Suck*}
- Goal test?? Check whether squares are clean.
- Path cost?? Number of actions to reach goal.

Example: 8-puzzle

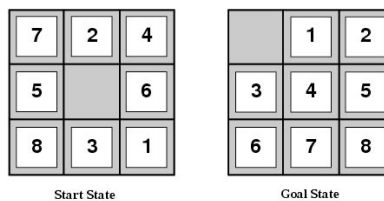| 7 | 2 | 4 |
| 5 | | 6 |
| 8 | 3 | 1 |

Start State

	1	2
3	4	5
6	7	8

Goal State

- States??
- Initial state??
- Actions??
- Goal test??
- Path cost??

Example: 8-puzzle

7	2	4
5		6
8	3	1

Start State
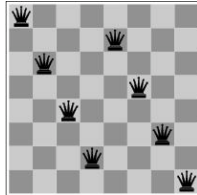
	1	2
3	4	5
6	7	8

Goal State

- States?? Integer location of each tile
- Initial state?? Any state can be initial
- Actions?? {*Left, Right, Up, Down*}
- Goal test?? Check whether goal configuration is reached
- Path cost?? Number of actions to reach goal

Example: 8-queens problem

- States??
- Initial state??
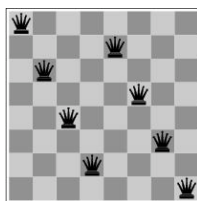- Actions??
- Goal test??
- Path cost??

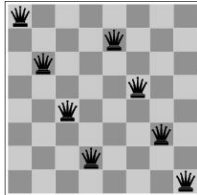# Example: 8-queens problem

Incremental formulation vs. complete-state formulation
- States??
- Initial state??
- Actions??
- Goal test??
- Path cost??

Example: 8-queens problem
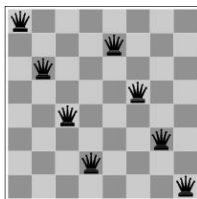

Incremental formulation

- States?? Any arrangement of 0 to 8 queens on the board
- Initial state?? No queens
- Actions?? Add queen in empty square
- Goal test?? 8 queens on board and none attacked
- Path cost?? None

  3×10^{14} possible sequences to investigate

Example: 8-queens problem
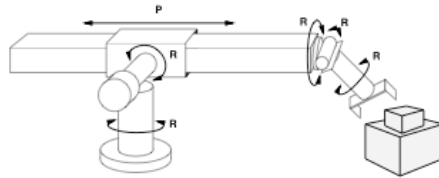

Incremental formulation (alternative)

- States?? n ($0 \leq n \leq 8$) queens on the board, one per column in the n leftmost columns with no queen attacking another.
- Actions?? Add queen in leftmost empty column such that is not attacking other queens

  2057 possible sequences to investigate; Yet makes no difference when n=100

Example: robot assembly
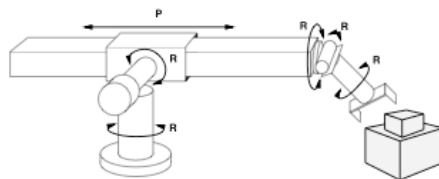

- States??
- Initial state??
- Actions??
- Goal test??
- Path cost??

Example: robot assembly

- States?? Real-valued coordinates of robot joint angles; parts of the object to be assembled.
- Initial state?? Any arm position and object configuration.
- Actions?? Continuous motion of robot joints
- Goal test?? Complete assembly (without robot)
- Path cost?? Time to execute

Basic search algorithms

- How do we find the solutions of previous problems?
  - **Search the state space (remember complexity of space depends on state representation)**

  - **Here: search through *explicit tree generation***
    - ROOT= initial state.
    - Nodes and leafs generated through successor function.

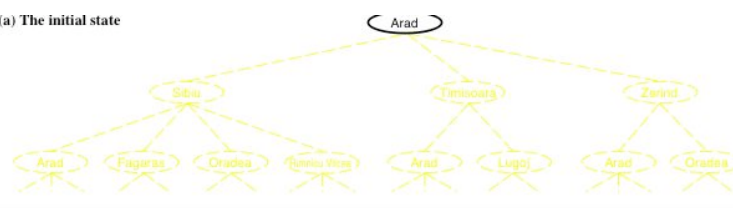 - **In general search generates a graph (same state through multiple paths)**

Simple tree search example

(a) The initial state

function TREE-SEARCH(*problem, strategy*) **return** a solution or failure
Initialize search tree to the *initial state* of the *problem*
do
        if no candidates for expansion **then return** *failure*
        choose leaf node for expansion according to *strategy*
        if node contains goal state **then return** *solution*
        else expand the node and add resulting nodes to the search tree
enddo

Simple tree search example

(b) After expanding Arad

function TREE-SEARCH(*problem, strategy*) **return** a solution or failure
    Initialize search tree to the *initial state* of the *problem*
    do
        if no candidates for expansion **then return** *failure*
        choose leaf node for expansion according to *strategy*
        if node contains goal state **then return** *solution*
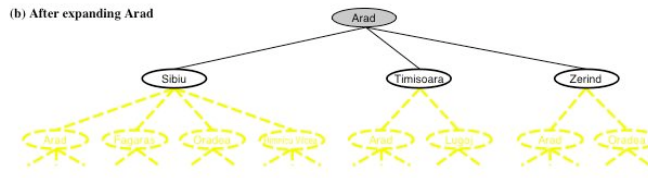   **else** expand the node and add resulting nodes to the search tree
    enddo

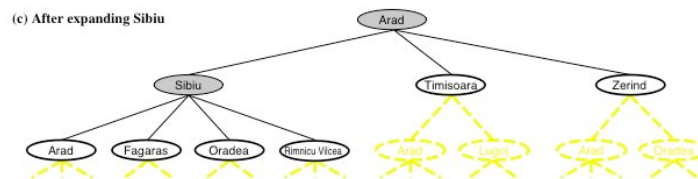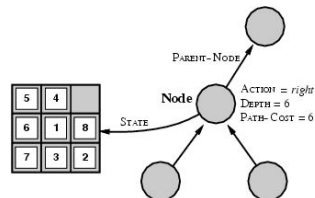# Simple tree search example

(c) After expanding Sibiu

function TREE-SEARCH(*problem, strategy*) **return** a solution or failure
    Initialize search tree to the *initial state* of the *problem*
    do
        if no candidates for expansion **then return** *failure*
        choose leaf node for expansion according to *strategy* ← **Determines search process!!**
        if node contains goal state **then return** *solution*
        else expand the node and add resulting nodes to the search tree
    enddo

State space vs. search tree

- A *state* is a (representation of) a physical configuration
- A *node* is a data structure belong to a search tree
  - A node has a parent, children, ... and ncludes path cost, depth, ...
  - Here *node= <state, parent-node, action, path-cost, depth>*
  - *FRINGE*= contains generated nodes which are not yet expanded.
    - White nodes with black outline

Tree search algorithm

function TREE-SEARCH(*problem,fringe*) **return** a solution or failure
   fringe ← INSERT(MAKE-NODE(INITIAL-STATE[*problem*]), *fringe*)
   loop do
      if EMPTY?(*fringe*) **then return** failure
      node ← REMOVE-FIRST(*fringe*)
      if GOAL-TEST[*problem*] applied to STATE[*node*] succeeds
            then return SOLUTION(*node*)
      fringe ← INSERT-ALL(EXPAND(*node*, *problem*), *fringe*)

Tree search algorithm (2)

function EXPAND(*node,problem*) **return** a set of nodes
    successors ← the empty set
    for each <*action*, *result*> **in** SUCCESSOR-FN[*problem*](STATE[*node*]) **do**
        s ← a new NODE
        STATE[*s*] ← *result*
        PARENT-NODE[*s*] ← *node*
        ACTION[*s*] ← *action*
        PATH-COST[*s*] ← PATH-COST[*node*] + STEP-COST(*node, action,s*)
        DEPTH[*s*] ← DEPTH[*node*]+1
        add *s* to *successors*
    return *successors*

Search strategies

- A strategy is defined by picking the order of node expansion.
- Problem-solving performance is measured in four ways:
  - □ **Completeness;** *Does it always find a solution if one exists?*
  - □ **Optimality;** *Does it always find the least-cost solution?*
  - □ **Time Complexity;** *Number of nodes generated/expanded?*
  - □ **Space Complexity;** *Number of nodes stored in memory during search?*
- Time and space complexity are measured in terms of problem difficulty defined by:
  - □ *b - maximum branching factor of the search tree*
  - □ *d - depth of the least-cost solution*
  - □ *m - maximum depth of the state space (may be ∞)*

Uninformed search strategies

- (a.k.a. blind search) = use only information available in problem definition.
  - □ **When strategies can determine whether one non-goal state is better than another → informed search.**
- Categories defined by expansion algorithm:
  - □ **Breadth-first search**
  - □ **Uniform-cost search**
  - □ **Depth-first search**
  - □ **Depth-limited search**
  - □ **Iterative deepening search.**
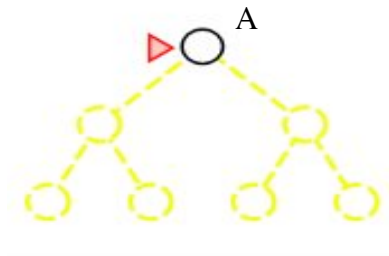 - □ **Bidirectional search**

BF-search, an example

- Expand *shallowest* unexpanded node
- Implementation: *fringe* is a FIFO queue

BF-search, an example
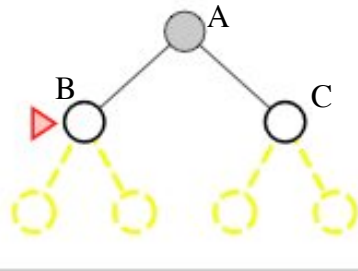
- Expand *shallowest* unexpanded node
- Implementation: *fringe* is a FIFO queue

BF-search, an example

- Expand *shallowest* unexpanded node
- Implementation: *fringe* is a FIFO queue

BF-search, an example

- Expand *shallowest* unexpanded node
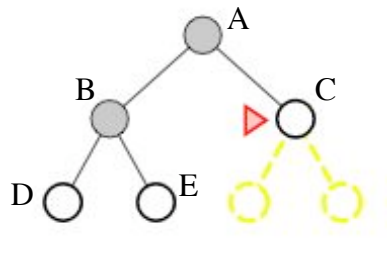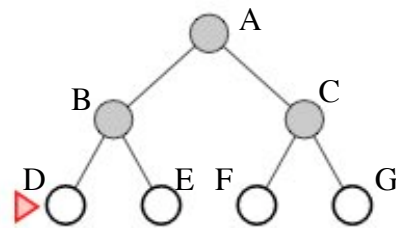- Implementation: *fringe* is a FIFO queue

BF-search; evaluation

- Completeness:
  - ☐ *Does it always find a solution if one exists?*
  - ☐ **YES**
    - If shallowest goal node is at some finite depth d
    - Condition: If *b* is finite
      - ☐ **(maximum num. Of succ. nodes is finite)**

BF-search; evaluation

- Completeness:
  - **YES (if *b* is finite)**
- Time complexity:
  - **Assume a state space where every state has *b* successors.**
    - root has *b* successors, each node at the next level has again b successors (total b^2), …
    - Assume solution is at depth *d*
    - Worst case; expand all but the last node at depth *d*
    - Total numb. of nodes generated:

$$b + b^2 + b^3 + ... + b^d + (b^{d+1} - b) = O(b^{d+1})$$

BF-search; evaluation

- Completeness:
  - **YES (if *b* is finite)**
- Time complexity:
  - **Total numb. of nodes generated:**

$$b + b^2 + b^3 + ... + b^d + (b^{d+1} - b) = O(b^{d+1})$$

- Space complexity:
  - **Idem if each node is retained in memory**

BF-search; evaluation

- Completeness:
  - **YES (if *b* is finite)**
- Time complexity:
  - **Total numb. of nodes generated:**

$$b + b^2 + b^3 + ... + b^d + (b^{d+1} - b) = O(b^{d+1})$$

- Space complexity:
  - **Idem if each node is retained in memory**
- Optimality:
  - ***Does it always find the least-cost solution?***
  - **In general YES**
    - unless actions have different cost.

BF-search; evaluation

- Two lessons:
  - **Memory requirements are a bigger problem than its execution time.**
  - **Exponential complexity search problems cannot be solved by uninformed search methods for any but the smallest instances.**

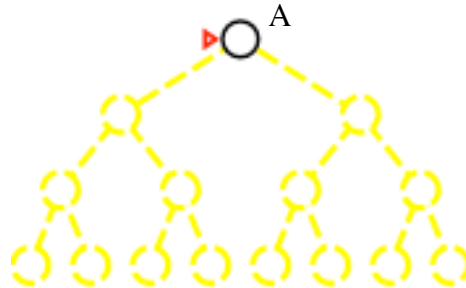| DEPTH2	NODES	TIME	MEMORY
2	1100	0.11 seconds	1 megabyte
4	111100	11 seconds	106 megabytes
6	10^7	19 minutes	10 gigabytes
8	10^9	31 hours	1 terabyte
10	10^{11}	129 days	101 terabytes
12	10^{13}	35 years	10 petabytes
14	10^{15}	3523 years	1 exabyte

Uniform-cost search

- Extension of BF-search:
  - **Expand node with *lowest path cost***

- Implementation: *fringe* = queue ordered by path cost.

- UC-search is the same as BF-search when all step-costs are equal.

Uniform-cost search

- Completeness:
  - **YES, if step-cost > ε (smal positive constant)**
- Time complexity:
  - **Assume C* the cost of the optimal solution.**
  - **Assume that every action costs at least ε**
  - **Worst-case:** $O(b^{C*/\varepsilon})$

- Space complexity:
  - **Idem to time complexity**
- Optimality:
  - **nodes expanded in order of increasing path cost.**
  - **YES, if complete.**

DF-search, an example

- Expand *deepest* unexpanded node
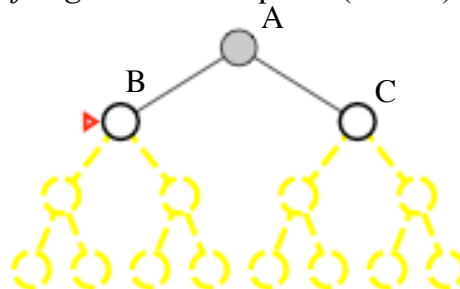- Implementation: *fringe* is a LIFO queue (=stack)

A

DF-search, an example

- Expand *deepest* unexpanded node
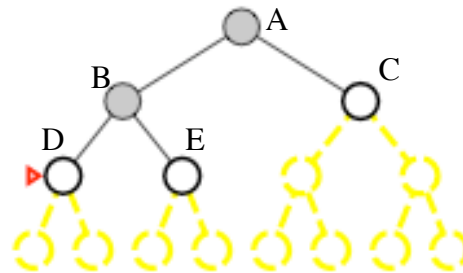- Implementation: *fringe* is a LIFO queue (=stack)

A

B   C

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)

A

B C

D E
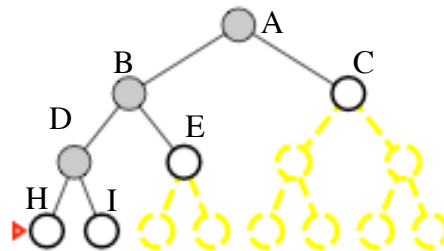
H I

TLo (IRIDIA)

DF-search, an example

- Expand *deepest* unexpanded node
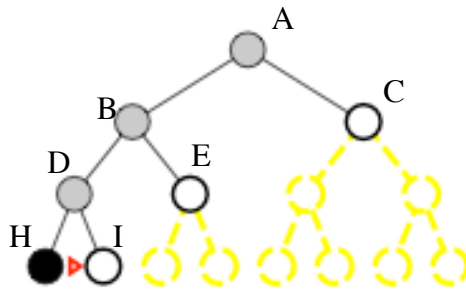- Implementation: *fringe* is a LIFO queue (=stack)

A

B C

D E

H I

TLo (IRIDIA)

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)
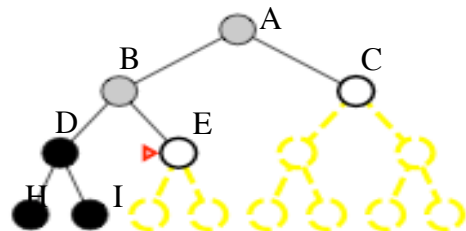
A
B
C
D
E
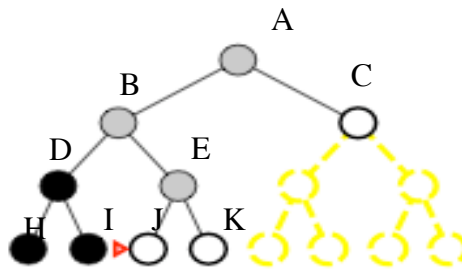H I J K

TLo (IRIDIA)

49

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)

A
B
C
D
E
H I J K

TLo (IRIDIA)

50

25

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)

TLo (IRIDIA)

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)

TLo (IRIDIA)

DF-search, an example

- Expand *deepest* unexpanded node
- Implementation: *fringe* is a LIFO queue (=stack)

TLo (IRIDIA)

DF-search, an example

- Expand *deepest* unexpanded node
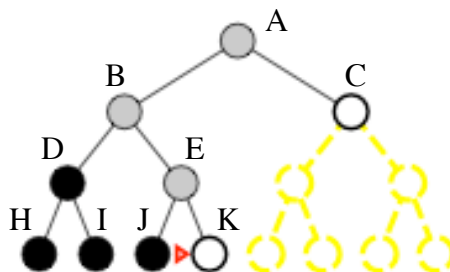- Implementation: *fringe* is a LIFO queue (=stack)

TLo (IRIDIA)
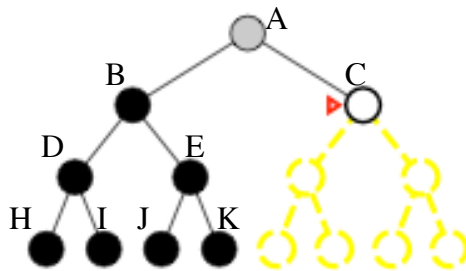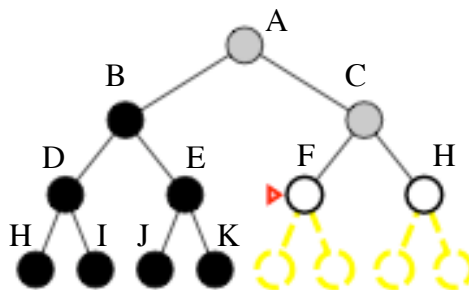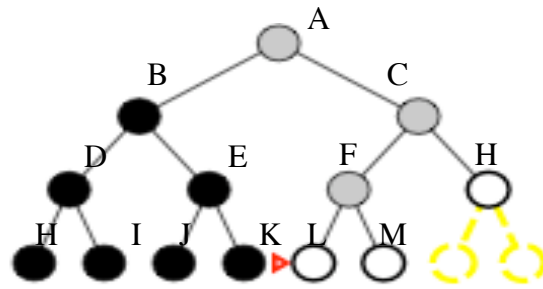
DF-search; evaluation

- Completeness;
  - *Does it always find a solution if one exists?*
  - **NO**
    - *unless* search space is finite and no loops are possible.

DF-search; evaluation

- Completeness;
  - **NO unless search space is finite.**
- Time complexity; $O(b^m)$
  - **Terrible if *m* is much larger than *d* (depth of optimal solution)**
  - **But if many solutions, then faster than BF-search**

DF-search; evaluation

- Completeness;
  - **NO unless search space is finite.**
- Time complexity; $O(b^m)$
- Space complexity; $O(bm + 1)$
  - **Backtracking search uses even less memory**
    - One successor instead of all *b*.

DF-search; evaluation

- Completeness;
  - **NO unless search space is finite.**
- Time complexity; $O(b^m)$
- Space complexity; $O(bm + 1)$
- Optimallity; No
  - **Same issues as completeness**
  - **Assume node J and C contain goal states**

Depth-limited search

- Is DF-search with depth limit l.
  - □ i.e. nodes at depth l have no successors.
  - □ Problem knowledge can be used
- Solves the infinite-path problem.
- If $l < d$ then incompleteness can result.
- If $l > d$ then not optimal.
- Time complexity: $O(b^l)$
- Space complexity: $O(bl)$

Depth-limited algorithm

function DEPTH-LIMITED-SEARCH(*problem,limit*) **return** a solution or failure/cutoff
    return RECURSIVE-DLS(MAKE-NODE(INITIAL-STATE[*problem*]),*problem,limit*)

function RECURSIVE-DLS(*node, problem, limit*) **return** a solution or failure/cutoff
   cutoff_occurred? ← false
   if GOAL-TEST[*problem*](STATE[*node*]) **then return** SOLUTION(*node*)
   else if DEPTH[*node*] == *limit* then return *cutoff*
   else for each *successor* **in** EXPAND(*node, problem*) **do**
       result ← RECURSIVE-DLS(*successor, problem, limit*)
      if *result* == *cutoff* **then** *cutoff_occurred?* ← true
      else if *result* ≠ *failure* **then return** *result*
   if *cutoff_occurred?* **then return** *cutoff* **else return** failure

Iterative deepening search

- ■ What?
  - □ **A general strategy to find best depth limit *l*.**
    - ■ Goals is found at depth d, the depth of the shallowest goal-node.
  - □ **Often used in combination with DF-search**
- ■ Combines benefits of DF- en BF-search

Iterative deepening search

function ITERATIVE_DEEPENING_SEARCH(*problem*) **return** a solution or failure

   inputs: *problem*

   for *depth* ← 0 to ∞ **do**
       result ← DEPTH-LIMITED_SEARCH(*problem*, *depth*)
       if *result* ≠ *cuttoff* **then return** *result*

ID-search, example

- Limit=0

TLo (IRIDIA)

ID-search, example

- Limit=1

TLo (IRIDIA)

ID-search, example

- Limit=2

TLo (IRIDIA)   **65**

ID-search, example

- Limit=3

TLo (IRIDIA)   **66**

33

ID search, evaluation

- Completeness:
  - **YES (no infinite paths)**

ID search, evaluation

- Completeness:
  - **YES (no infinite paths)**
- Time complexity:
  - **Algorithm seems costly due to repeated generation of certain states.**
  - **Node generation:** $O(b^d)$
    - level d: once
    - level d-1: 2
    - level d-2: 3
    - …
    - level 2: d-1
    - level 1: d

$$N(IDS) = (d)b + (d-1)b^2 + ... + (1)b^d$$

$$N(BFS) = b + b^2 + ... + b^d + (b^{d+1} - b)$$

Num. Comparison for b=10 and d=5 solution at far right

$$N(IDS) = 50 + 400 + 3000 + 20000 + 100000 = 123450$$

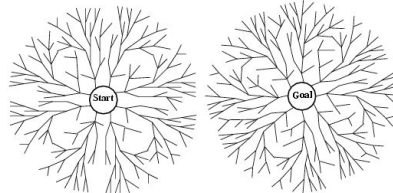$$N(BFS) = 10 + 100 + 1000 + 10000 + 100000 + 999990 = 1111100$$

ID search, evaluation

- Completeness:
    - **YES (no infinite paths)**
- Time complexity: $O(b^d)$
- Space complexity: $O(bd)$
    - **Cfr. depth-first search**

ID search, evaluation

- Completeness:
    - **YES (no infinite paths)**
- Time complexity: $O(b^d)$
- Space complexity: $O(bd)$
- Optimality:
    - **YES if step cost is 1.**
    - **Can be extended to iterative lengthening search**
        - Same idea as uniform-cost search
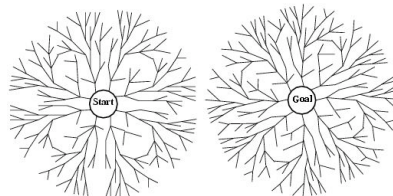   - Increases overhead.

Bidirectional search

- Two simultaneous searches from start an goal.
  - **Motivation:** $b^{d/2} + b^{d/2} \neq b^d$
- Check whether the node belongs to the other fringe before expansion.
- Space complexity is the most significant weakness.
- Complete and optimal if both searches are BF.

How to search backwards?

- The predecessor of each node should be efficiently computable.
  - **When actions are easily reversible.**

Summary of algorithms

Criterion	Breadth-First	Uniform-cost	Depth-First	Depth-limited	Iterative deepening	Bidirectional search
Complete?	YES*	YES*	NO	YES, if l ≥ d	YES	YES*
Time	b^{d+1}	$b^{C^*/e}$	b^m	b^l	b^d	$b^{d/2}$
Space	b^{d+1}	$b^{C^*/e}$	bm	bl	bd	$b^{d/2}$
Optimal?	YES*	YES*	NO	NO	YES	YES

Repeated states

- Failure to detect repeated states can turn a solvable problems into unsolvable ones.

Graph search algorithm

- Closed list stores all expanded nodes

function GRAPH-SEARCH(*problem,fringe*) **return** a solution or failure
    closed ← an empty set
    fringe ← INSERT(MAKE-NODE(INITIAL-STATE[*problem*]), *fringe*)
    loop do
        if EMPTY?(*fringe*) **then return** failure
        node ← REMOVE-FIRST(*fringe*)
        if GOAL-TEST[*problem*] applied to STATE[*node*] succeeds
                then return SOLUTION(*node*)
        if STATE[*node*] is not in *closed* **then**
            add STATE[*node*] to *closed*
            fringe ← INSERT-ALL(EXPAND(*node*, *problem*), *fringe*)

Graph search, evaluation

- Optimality:
  - **GRAPH-SEARCH discard newly discovered paths.**
    - This may result in a sub-optimal solution
    - YET: when uniform-cost search or BF-search with constant step cost
- Time and space complexity,
  - **proportional to the size of the state space**
    (may be much smaller than $O(b^d)$).
  - **DF- and ID-search with closed list no longer has linear space requirements since all nodes are stored in closed list!!**

Search with partial information

- Previous assumption:
  - □ **Environment is fully observable**
  - □ **Environment is deterministic**
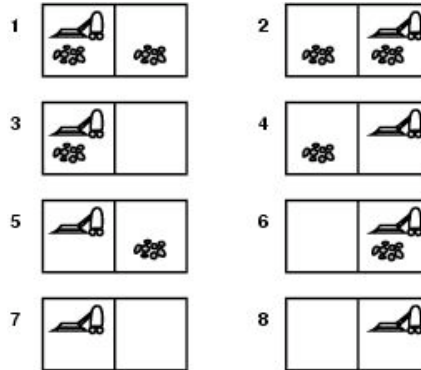 - □ **Agent knows the effects of its actions**

What if knowledge of states or actions is incomplete?

Search with partial information

- (SLIDE 7) Partial knowledge of states and actions:
  - □ *sensorless or conformant problem*
    - Agent may have no idea where it is; solution (if any) is a sequence.
  - □ *contingency problem*
    - Percepts provide *new* information about current state; solution is a tree or policy; often interleave search and execution.
    - If uncertainty is caused by actions of another agent: *adversarial problem*
  - □ *exploration problem*
    - When states and actions of the environment are unknown.

Conformant problems

- start in {1,2,3,4,5,6,7,8} e.g Right goes to {2,4,6,8}. Solution??
  - *[Right, Suck, Left,Suck]*
- *When the world is not fully observable: reason about a set of states that migth be reached*
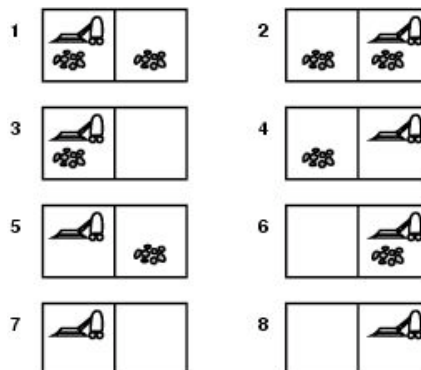 **=belief state**

Conformant problems

- Search space of belief states
- Solution = belief state with all members goal states.
- If S states then 2^S belief states.
- Murphy's law:
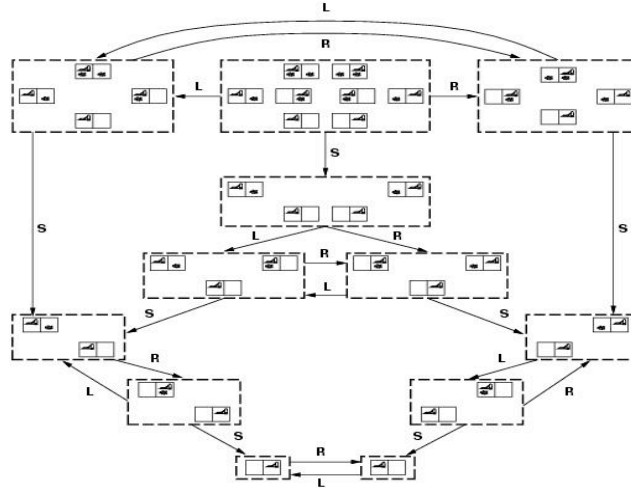 - *Suck can dirty a clear square.*

Belief state of vacuum-world
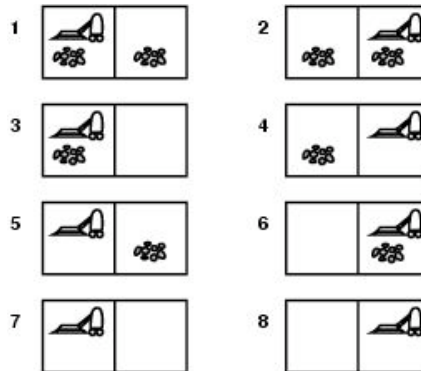

Contingency problems

- Contingency, start in {1,3}.
- Murphy's law, Suck *can* dirty a clean carpet.
- Local sensing: dirt, location only.
  - Percept = [L,Dirty] ={1,3}
  - [Suck] = {5,7}
  - [Right] ={6,8}
  - [Suck] in {6}={8} (Success)
  - BUT [Suck] in {8} = failure
- Solution??
  - Belief-state: no fixed action sequence guarantees solution
- Relax requirement:
  - [*Suck, Right,* if *[R,dirty]* then *Suck*]
  - Select actions based on contingencies arising during execution.